## Lecture 2: <br> LanaOa Normal-form game (Strategic-form game) with pure strategies <br> Yi, Yung (이융) <br> KAIST, Electrical Engineering http://lanada.kaist.ac.kr <br> yiyung@kaist.edu

## Key Words

- Normal-form (Strategic form) Game
- Matrix game
- Strategy spaces are discrete
- Continuous-kernel game
- Strategy spaces are continuous
- Strictly dominated strategies
- Pure/Mixed strategy
- Saddle point, Nash equilibrium

Matrix Game: Pure Strategy

## Matrix Game

- Representation of a game
- Simultaneous play
- players analyze the game and write their strategy on a paper
- Combination of strategies determines payoff



## More Formal Game Definition

- Normal form (strategic) game
- a finite set $\boldsymbol{N}$ of players
- a set strategies $A_{i}$ for each player i
- payoff function $u_{i}(s)$ for each player $i \in N$
- where $s \in A=\times_{j \in N} A_{j}$ is the set of strategies chosen by all players

$$
i \in N
$$

- $A$ is the set of all possible outcomes
- $s \in A$ is a set of strategies chosen by players
- defines an outcome
- $u_{i}: A \rightarrow \mathfrak{R}$


## Two-person Zero-sum Games

- One of the first games studied
- most well understood type of game
- Players interest are strictly opposed
- what one player gains what the other loses
- game matrix has single entry (gain to player 1)
- Intuitive solution concept
- players maximize gains
- unique solution


## Solution Concept

- A formal rule for predicting how a game will be played
- Describes which strategies will be adopted by palyers, and thus the result of the game
- Many kinds of solution concepts
- People's perspectives are different
- It does not talk about how players reach a solution concept
- Thus, naturally, it is an "equilibrium concept".


## Analyzing the Game: Domination

- Player 1 maximizes matrix entry, while player 2 minimizes
 (dominated by B)


## Dominance

- Strategy S strictly dominates a strategy T
- if every possible outcome when $S$ is chosen is better than the corresponding outcome when T is chosen
- Dominance Principle
- rational players never choose strictly dominated strategies
- Idea: Solve the game by eliminating strictly dominated strategies!
- iterated removal


## Solving the Game

- Iterated removal of strictly dominated strategies

- Player 1 cannot remove any strategy (neither T or B dominates the other)
- Player 2 can remove strategy $R$ (dominated by $M$ )
- Player 1 can remove strategy $T$ (dominated by $B$ )
- Player 2 can remove strategy $L$ (dominated by $M$ )
- Solution: $P_{1} \rightarrow B, P_{2} \rightarrow M$
- payoff of 2


## Solving the Game

- Removal of strictly dominates strategies does not always work
- Consider the game

|  | Player 2 |  |  |  |
| :---: | :---: | :---: | :---: | :---: |
|  | A | B | D |  |
| A | 12 | -1 | 0 |  |
| C | 5 | 2 | 3 |  |
| D | -16 | 0 | -1 |  |

- Neither player has dominated strategies
- Requires another solution concept


## Analyzing the Game

## Player 2



Outcome ( $C, B$ ) seems
"stable"

- saddle point of game


## Saddle Points

- An outcome is a saddle point
- if it is both less than or equal to any value in its row and greater than or equal to any value in its column
- Saddle Point Principle
- Players should choose outcomes that are saddle points of the game
- Value of the game
- value of saddle point outcome if it exists


## Why Play Saddle Points?



- If player 1 believes player 2 will play B
- player 1 should play best response to $B$ (which is $C$ )
- If player 2 believes player 1 will play C
- player 2 should play best response to C (which is B )


## Why Play Saddle Points?



- Why should player 1 believe player 2 will play $B$ ?
- playing B guarantees player 2 loses at most $v$ (which is 2 )
- Why should player 2 believe player 1 will play C?
- playing C guarantees player 1 wins at least $v$ (which is 2 )


## Powerful arguments to play saddle point!

## Solving the Game (min-max algorithm)

## Player 2

|  | A | B | $C$ | $D$ |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| Player 1 | 4 | 3 | 2 | 5 | 2 |
| B | -10 | 2 | 0 | -1 | -10 |
| $C$ | 7 | 5 | 1 | 3 | 1 |
| $D$ | 0 | 8 | -4 | -5 | -5 |
|  | 7 | 8 | 2 | 5 |  |

- choose maximum entry in each column
- choose the minimum among these
- this is the minimax value
- if minimax == maximin, then this is the saddle point of game


## Multiple Saddle Points

- In general, game can have multiple saddle points



## Games With no Saddle Points



## Two-person Non-zero Sum Games

- Players are not strictly opposed
- payoff sum is non-zero

$$
\text { Player } 2
$$

|  |  | $A$ | $B$ |
| :---: | :---: | :---: | :---: |
| Player 1 | A | 3,4 | 2,0 |
|  | $B$ | 5,1 | $-1,2$ |

- Situations where interest is not directly opposed


## What is the Solution?

- Ideas of zero-sum game: saddle points
- pure strategy equilibrium
- no pure strategy eq.



## Nash equilibrium

- A Nash equilibrium is a strategy profile s* with the property that no player $i$ can do better by choosing a strategy different from s*, given that every other player $j \neq i$.
- In other words, for each player $i$ with payoff function $u_{i}$,

$$
u_{i}\left(s_{i}^{*}, s_{-i}^{*}\right) \geq u_{i}\left(s_{i}, s_{-i}^{*}\right), \forall s_{i} \in \mathcal{S}_{i}
$$

- No user can change its payoff by unilaterally changing its strategy, i.e., changing its strategy while $\mathbf{s}_{-\mathrm{i}}$ is fixed


## Multiple Solution Problem

- Games can have multiple equilibria
- not equivalent:
- payoff is different
- not interchangeable:
- playing an equilibrium strategy does not lead to equilibrium



## Ex 1: Coordination game

- Two drivers, driving towards each other

|  | Left | Right |
| :---: | :---: | :---: |
| Left | 1,1 | 0,0 |
| Right | 0,0 | 1,1 |
|  |  |  |

## Ex 2: Matching Pennies game

- Each player shows her coin.
- Same side $\rightarrow$ Player 1 pockets both, and Player 2 does otherwi se.

|  | Heads | Tails |
| :---: | :---: | :---: |
| Heads | $1,-1$ | $-1,1$ |
| Tails | $-1,1$ | $1,-1$ |
|  |  |  |

## Ex 3: Battle of the Sexes Game

- Tries to see a movie
- Husband: "Lethal Weapon", Wife: "Wondrous Love"

Husband


## Summary

