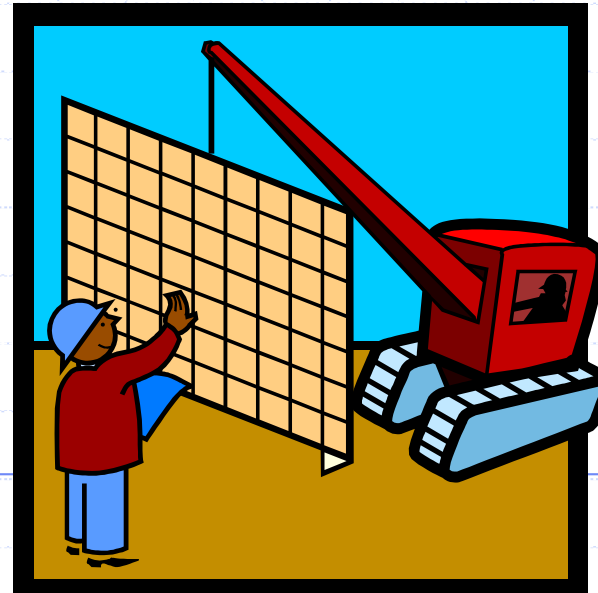


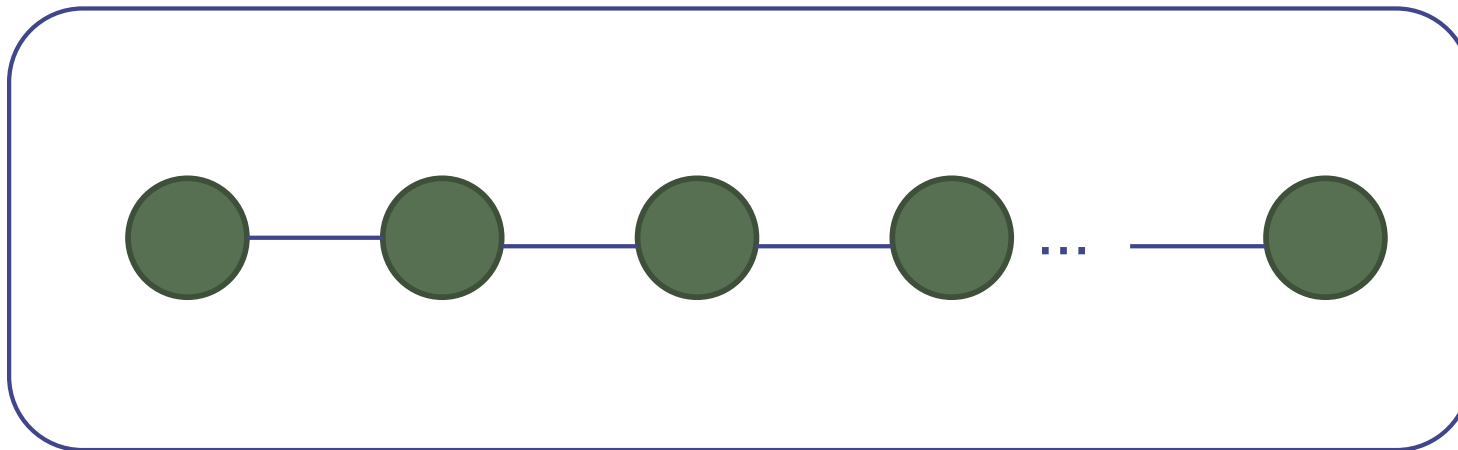
# Vector, List and Sequence



# Overview and Reading

---

- ◆ Reading: Chapters: 6.1, 6.2, and 6.3
- ◆ A data structure that stores  $n$  elements in a linear order
  - Called list or sequence
- ◆ Didn't we learn "array" and "linked list"?
  - We are talking about more abstract ADTs than them



# Three ADTs

---

## ◆ Vector (also called Array List)

- Access each element using a notion of **index in  $[0, n-1]$**
- Index of element  $e$ : the number of elements that are before  $e$
- Typically we use the “**index**” (e.g., [ ])
- A more general ADT than “array”

## ◆ List

- Not using an index to access, but use a node to access
- Insert a new element  $e$  before some “**position**”  $p$
- A more general ADT than “linked list”

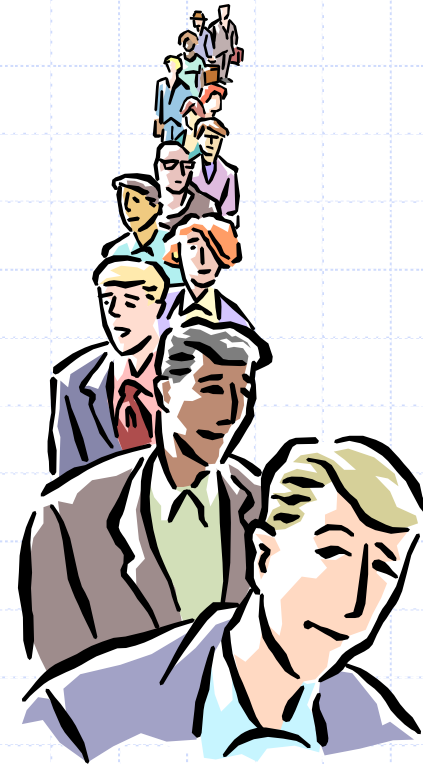
## ◆ Sequence

- Can access an element as vector and list (using both **index** and **position**)

## ◆ (Note) Can implement the above ADTs using various ways

- array, singly linked list, doubly linked list, circular linked list

# Vectors (or Array Lists)



# The Array List ADT

---

## □ The **Vector** or **Array List**

ADT extends the notion of array by storing a sequence of objects

□ An element can be accessed, inserted or removed by specifying its **index** (number of elements preceding it)

□ An exception is thrown if an incorrect index is given (e.g., a negative index)

## ◆ Main methods:

- **at**(integer i): returns the element at index i without removing it
- **set**(integer i, object o): replace the element at index i with o
- **insert**(integer i, object o): insert a new element o to have index i
- **erase**(integer i): removes element at index i

## ◆ Additional methods:

- **size**()
- **empty**()

# Applications of Array Lists

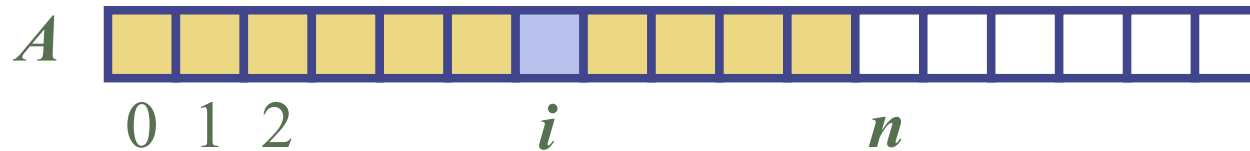
---

- ◆ Direct applications
  - Sorted collection of objects (elementary database)
- ◆ Indirect applications
  - Auxiliary data structure for algorithms
  - Component of other data structures
- ◆ Basically, every place where you can use “array”.

# Array-based Implementation of Vector

---

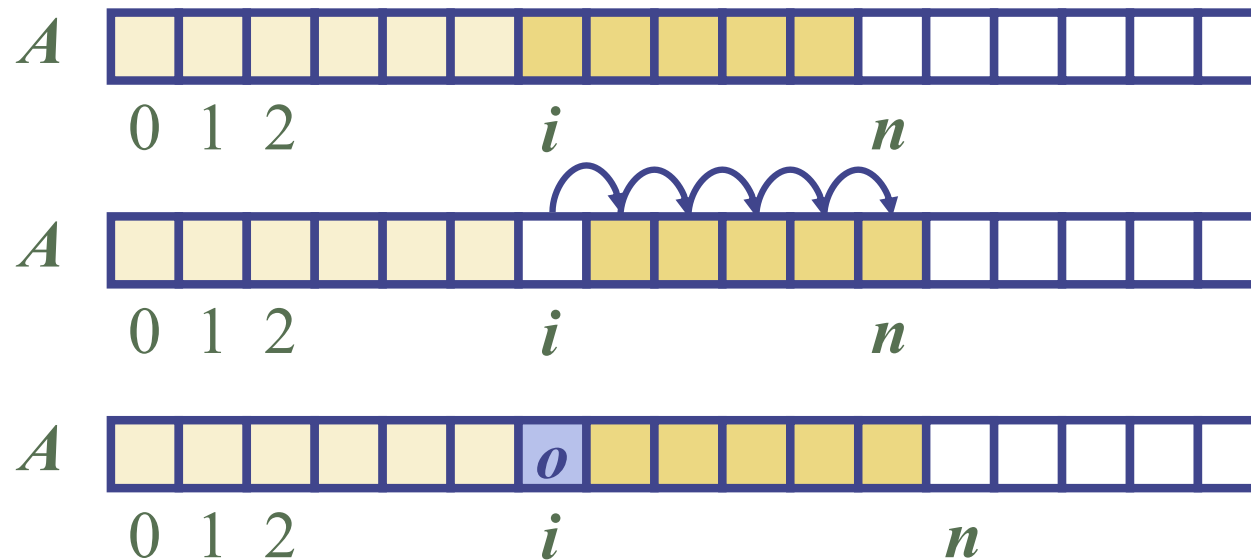
- ◆ Use an array  $A$  of size  $N$
- ◆ A variable  $n$  keeps track of the size of the array list (number of elements stored)
- ◆ Operation  $at(i)$  is implemented in  $O(1)$  time by returning  $A[i]$
- ◆ Operation  $set(i,o)$  is implemented in  $O(1)$  time by performing  $A[i] = o$



# Insertion

---

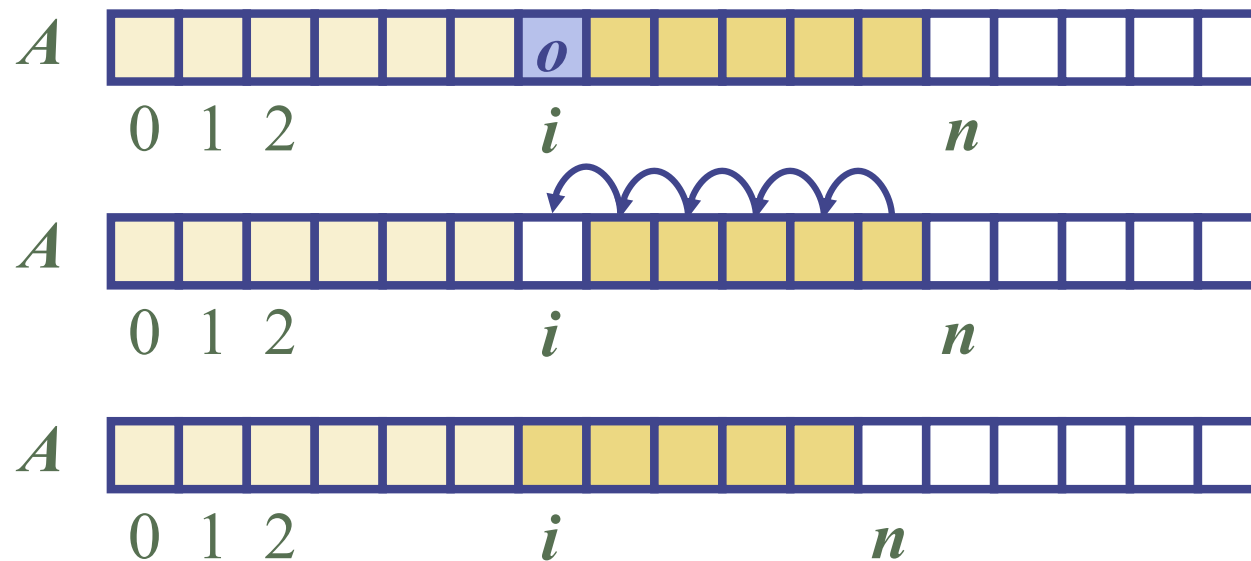
- ◆ In operation *insert*( $i, o$ ), we need to make room for the new element by shifting forward the  $n - i$  elements  $A[i], \dots, A[n - 1]$
- ◆ In the worst case ( $i = 0$ ), this takes  $O(n)$  time





# Element Removal

- ◆ In operation *erase*( $i$ ), we need to fill the hole left by the removed element by shifting backward the  $n - i - 1$  elements  $A[i + 1], \dots, A[n - 1]$
- ◆ In the worst case ( $i = 0$ ), this takes  $O(n)$  time



# Performance

---

- ◆ In the array-based implementation of an array list:
  - The space used by the data structure is  $O(n)$
  - *size*, *empty*, *at* and *set* run in  $O(1)$  time
  - *insert* and *erase* run in  $O(n)$  time in worst case
- ◆ If we use the array in a circular fashion, operations *insert*(0,  $x$ ) and *erase*(0,  $x$ ) run in  $O(1)$  time
- ◆ In an *insert* operation, when the array is full, instead of throwing an exception, we can replace the array with a larger one

# Growable Array-based Array List

- In an **insert(o)** operation (without an index), we always insert at the end
- When the array is full, we replace the array with a larger one
- How large should the new array be?
  - **Incremental strategy:** increase the size by a constant  $c$
  - **Doubling strategy:** double the size

```
Algorithm insert(o)  
  if  $t = S.length - 1$  then  
     $A \leftarrow$  new array of  
      size ...  
    for  $i \leftarrow 0$  to  $n-1$  do  
       $A[i] \leftarrow S[i]$   
     $S \leftarrow A$   
     $n \leftarrow n + 1$   
     $S[n-1] \leftarrow o$ 
```

- ◆ For size  $n$  array, “re-grow” operation requires  $n$  copies

# Which is better?: Incremental or Doubling

---

## ◆ Comparison Method 1

- Given the current size of  $S = n$
- Worst-case running time
  - ◆ Incremental strategy:  $O(1)$
  - ◆ Doubling strategy:  $O(n)$

## ◆ Are you happy?

- Happy if your focus is really the worst-case
- Unhappy
  - ◆ For doubling strategy, the total number of resizing array size would be small

## ◆ Can we reconsider the analysis method?

# Which is better?: Incremental or Doubling

---

## ◆ Comparison Method2

- Compute the total time  $T(n)$  needed to perform a series of  $n$  insert(o) operations
- Assume that we start with an empty stack represented by an array of size 1

## ◆ We call amortized time of an insert operation **the average time taken by an insert over the series of operations**, i.e., $T(n)/n$

- This can be a fairer comparison in some cases

## ◆ Amortized analysis (분할상환분석 in Wiki)

# Incremental Strategy Analysis

---

- ◆ We replace the old array with a new one  $k = n/c$  times
- ◆ The total time  $T(n)$  of a series of  $n$  insert operations is proportional to

$$\begin{aligned}n + c + 2c + 3c + 4c + \dots + kc &= \\n + c(1 + 2 + 3 + \dots + k) &= \\n + ck(k + 1)/2 &\end{aligned}$$

- ◆ Since  $c$  is a constant,  $T(n)$  is  $O(n + k^2)$ , i.e.,  $O(n^2)$
- ◆ The amortized time of an insert operation is  $O(n)$

# Doubling Strategy Analysis

---

- ◆ We replace the old array with a new one  $k = \log_2 n$  times

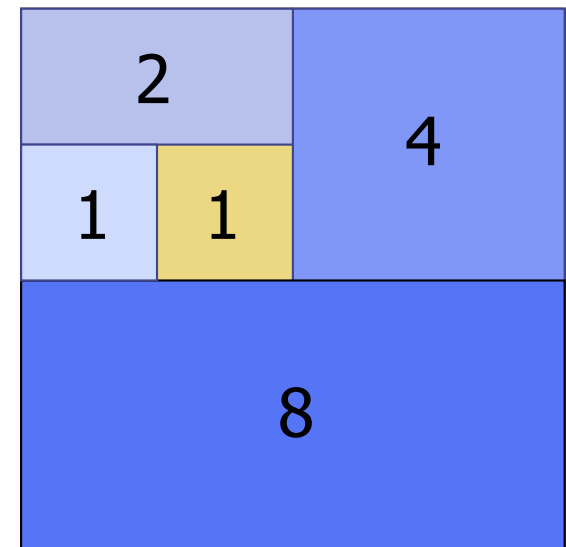
- ◆ The total time  $T(n)$  of a series of  $n$  insert operations is proportional to

$$\begin{aligned} n + 1 + 2 + 4 + 8 + \dots + 2^k &= \\ n + 2^{k+1} - 1 &= \\ 3n - 1 \end{aligned}$$

- ◆  $T(n)$  is  $O(n)$

- ◆ The amortized time of an insert operation is  $O(1)$

geometric series



# Professor, I have a question

---

- ◆ In “computing spans”, why didn’t you do amortized analysis?
- ◆ Can we do it?
- ◆ Is it meaningful?
  
- ◆ Think about this!
  - I am ready to discuss if you get your version of answer ready.

```
Algorithm spans2(X, n)           #
  S ← new array of n integers     n
  A ← new empty stack              1
  for i ← 0 to n – 1 do           n
    while ( $\neg A.empty()$  ∧
            $X[A.top()] \leq X[i]$ ) do n
      A.pop()                     n
    if A.empty() then             n
      S[i] ← i + 1                n
    else
      S[i] ← i – A.top()        n
      A.push(i)                   n
  return S                          1
```



# Vectors in C++ STL

---

```
#include <vector>           // provides definition of vector
using std::vector;        // make vector accessible

vector<int> myVector(100); // a vector with 100 integers
```

**vector(*n*):** Construct a vector with space for *n* elements; if no argument is given, create an empty vector.

**size():** Return the number of elements in *V*.

**empty():** Return true if *V* is empty and false otherwise.

**resize(*n*):** Resize *V*, so that it has space for *n* elements.

**reserve(*n*):** Request that the allocated storage space be large enough to hold *n* elements.

**operator[*i*]:** Return a reference to the *i*th element of *V*.

**at(*i*):** Same as *V*[*i*], but throw an `out_of_range` exception if *i* is out of bounds, that is, if  $i < 0$  or  $i \geq V.size()$ .

**front():** Return a reference to the first element of *V*.

**back():** Return a reference to the last element of *V*.

**push\_back(*e*):** Append a copy of the element *e* to the end of *V*, thus increasing its size by one.

**pop\_back():** Remove the last element of *V*, thus reducing its size by one.

Difference between  
resize() and reserve()?

# Logistics

---

- ◆ First programming assignment
  - Deadline: Sep, 19<sup>th</sup>
  
- ◆ Problem Solving Homework
  - Deadline: Oct, 1<sup>st</sup>
  
- ◆ You should keep reading the textbook
  
- ◆ Sep 24<sup>th</sup>, 26<sup>th</sup> : No class
  - Thanksgiving

# Last Class

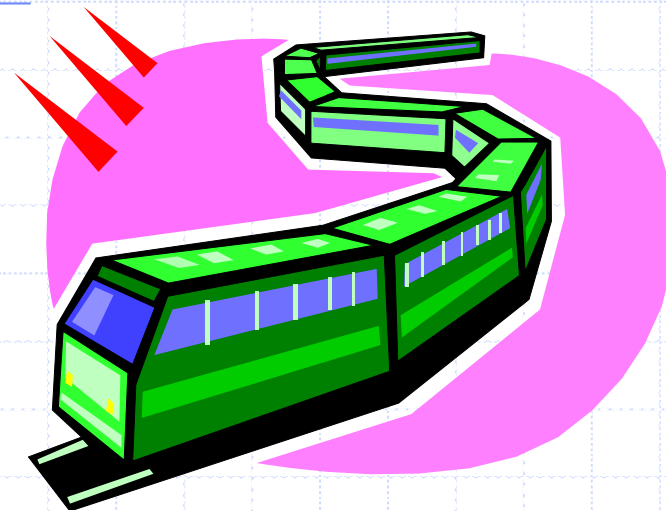
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## ◆ Vector and List

### ◆ Vector

- Access elements by “index”
- Incremental vs. Doubling Strategy
  - ◆ Amortized analysis

# Lists



# (Node) List ADT

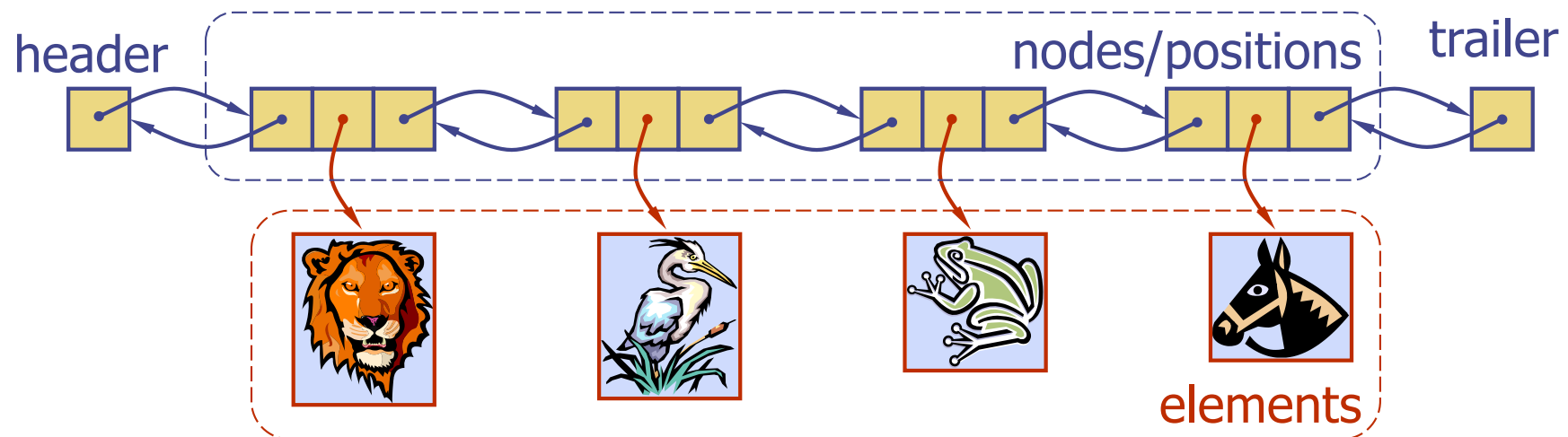
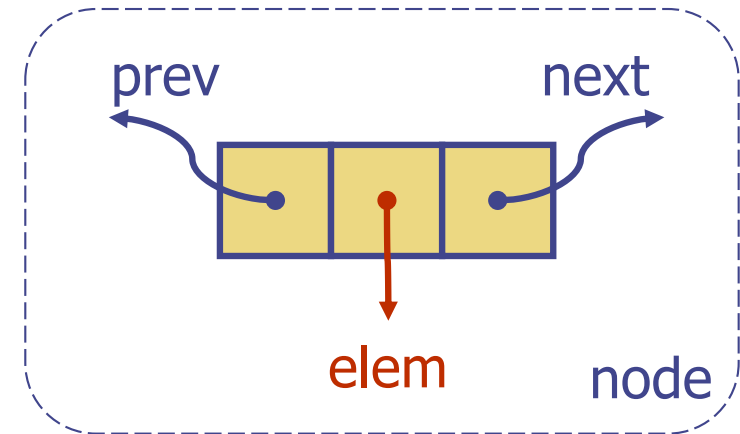
---

- ◆ The **Node List** ADT models a sequence of positions storing arbitrary objects
- ◆ It establishes a before/after relation between positions
- ◆ Generic methods:
  - **size()**, **empty()**
- Iterators:
  - begin(), end()
- Update methods:
  - **insertFront(e)**, **insertBack(e)**
  - **removeFront()**, **removeBack()**
- Iterator-based update:
  - insert(p, e)
  - remove(p)

(Question) No method for accessing a specific node?  
We will talk about this later

# Implementation based on DLL (covered this)

- ◆ A doubly linked list provides a natural implementation of the Node List ADT
- ◆ Nodes implement Position and store:
  - element
  - link to the previous node
  - link to the next node
- ◆ Special trailer and header nodes



# Performance

---

- ◆ In the implementation of the List ADT by means of a doubly linked list
  - The space used by a list with  $n$  elements is  $O(n)$
  - The space used by each position of the list is  $O(1)$
  - All the operations of the List ADT run in  $O(1)$  time

# Lists in C++ STL

---

```
#include <list>
using std::list;           // make list accessible
list<float> myList;       // an empty list of floats
```

`list(n)`: Construct a list with *n* elements; if no argument list is given, an empty list is created.

`size()`: Return the number of elements in *L*.

`empty()`: Return true if *L* is empty and false otherwise.

`front()`: Return a reference to the first element of *L*.

`back()`: Return a reference to the last element of *L*.

`push_front(e)`: Insert a copy of *e* at the beginning of *L*.

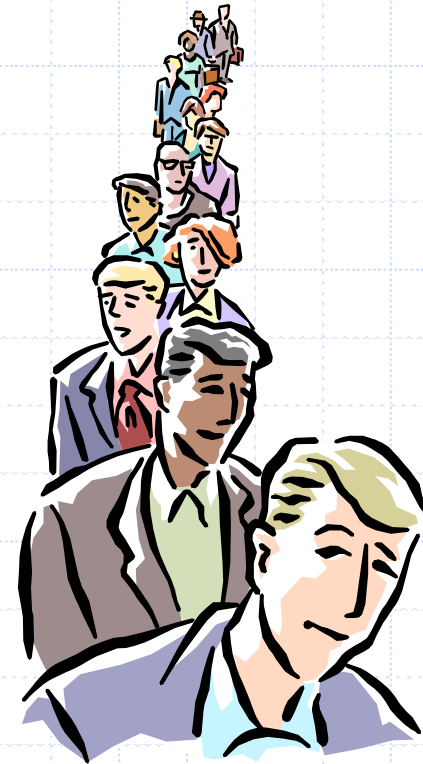
`push_back(e)`: Insert a copy of *e* at the end of *L*.

`pop_front()`: Remove the first element of *L*.

`pop_back()`: Remove the last element of *L*.



# Containers, Iterators, and Generic algorithms



# Sorting: Vector and List

---

◆ I want to find “yiyung” in Vector or List objects

```
vector<string> V(100);  
list<string> L(100);  
// some data insertion to V and L
```

```
//Design 1: different function  
find_vector(&V);  
find_list(&L);
```

```
//Design 2: function overloading  
find(&V);  
find(&L);
```

Do you like these? Why? Why not?

# This is how we can do in C++

---

```
#include <iostream>
#include <vector>
#include <string>
#include <algorithm>

using namespace::std;

int main()
{
    vector<string> vec_str;
    vec_str.push_back("is");
    vec_str.push_back("of");
    vec_str.push_back("the");
    vec_str.push_back("hello");

    vector<string>::iterator it;

    it =
        find(vec_str.begin(), vec_str.end(), "the");
    cout << "Print: " << *it << endl;

    it++;
    cout << "Print: " << *it << endl;

    return 0;
}
```

```
#include <iostream>
#include <list>
#include <string>
#include <algorithm>

using namespace::std;

int main()
{
    list<string> list_str;
    list_str.push_back("is");
    list_str.push_back("of");
    list_str.push_back("the");
    list_str.push_back("hello");

    list<string>::iterator it;

    it =
        find(list_str.begin(), list_str.end(), "the");
    cout << "Print: " << *it << endl;

    it++;
    cout << "Print: " << *it << endl;

    return 0;
}
```

It is cool. But why is it cool?

# Mysterious things

```
#include <iostream>
#include <vector>
#include <string>
#include <algorithm>

using namespace::std;

int main()
{
    vector<string> vec_str;
    vec_str.push_back("is");
    vec_str.push_back("of");
    vec_str.push_back("the");
    vec_str.push_back("hello");

    vector<string>::iterator it;

    it =
        find(vec_str.begin(), vec_str.end(), "the");
    cout << "Print: " << *it << endl;

    it++;
    cout << "Print: " << *it << endl;

    return 0;
}
```

```
#include <iostream>
#include <list>
#include <string>
#include <algorithm>

using namespace::std;

int main()
{
    list<string> list_str;
    list_str.push_back("is");
    list_str.push_back("of");
    list_str.push_back("the");
    list_str.push_back("hello");

    list<string>::iterator it;

    it =
        find(list_str.begin(), list_str.end(), "the");
    cout << "Print: " << *it << endl;

    it++;
    cout << "Print: " << *it << endl;

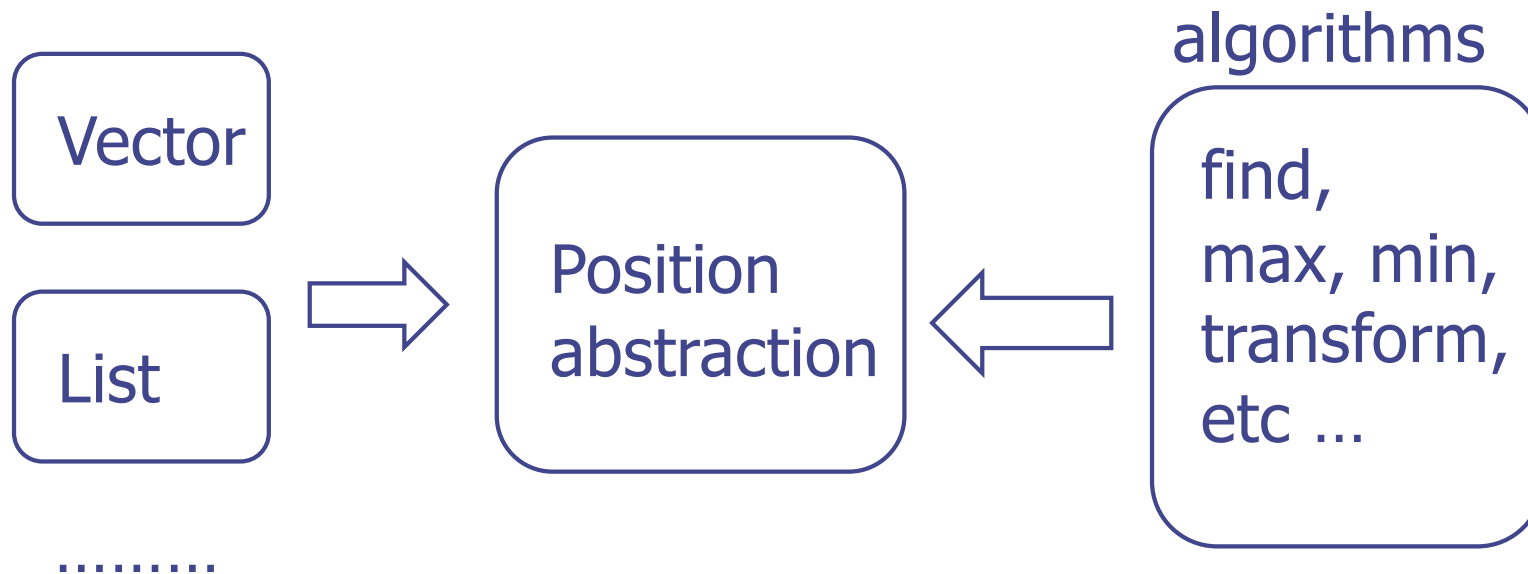
    return 0;
}
```

iterator? Looks like a "position" of vector or list. Hmm.....

# Goal and Design Challenge

---

- ◆ Lots of data structures (or classes in C++) that can contain various types of elements
  - “Container”
  - Examples: Vector, List, deque, set, map, etc ...



- ◆ How are you going to design this concept?
  - Again, from C++ STL designer’s perspective

# Position ADT

---

- The **Position** ADT models the notion of place within a data structure where a single object is stored
- It gives a unified view of diverse ways of storing data, such as
  - a cell of an array
  - a node of a linked list
- “A” method of accessing the element at position **p**:
  - object **p.element()**: returns the element at position
  - In C++ it is convenient to implement this as **\*p**
    - **Operator overloading**
- ◆ Implemented as “iterator” in C++

# Containers and Iterators in C++

---

- ◆ An **iterator** abstracts the process of scanning through a collection of elements
- ◆ A **container** is an abstract data structure that supports element access through iterators
  - Data structures that support iterators
  - Examples include Stack, Queue, Vector, List
  - **begin()**: returns an iterator to the first element
  - **end()**: return an iterator to an imaginary position just after the last element
- ◆ An iterator behaves like a pointer to an element
  - **\*p**: returns the element referenced by this iterator
  - **++p**: advances to the next element
- ◆ Extends the concept of **position** by adding a traversal capability

# Example codes again

---

```
#include <iostream>
#include <vector>
#include <string>
#include <algorithm>

using namespace::std;

int main()
{
    vector<string> vec_str;
    vec_str.push_back("is");
    vec_str.push_back("of");
    vec_str.push_back("the");
    vec_str.push_back("hello");

    vector<string>::iterator it;

    it =
        find(vec_str.begin(), vec_str.end(), "the");
    cout << "Print: " << *it << endl;

    it++;
    cout << "Print: " << *it << endl;

    return 0;
}
```

```
#include <iostream>
#include <list>
#include <string>
#include <algorithm>

using namespace::std;

int main()
{
    list<string> list_str;
    list_str.push_back("is");
    list_str.push_back("of");
    list_str.push_back("the");
    list_str.push_back("hello");

    list<string>::iterator it;

    it =
        find(list_str.begin(), list_str.end(), "the");
    cout << "Print: " << *it << endl;

    it++;
    cout << "Print: " << *it << endl;

    return 0;
}
```

Ah-ha, it's an iterator!



# Various Iterators

---

- ◆ (standard) iterator: allows read-write access to elements
- ◆ const iterator: provides read-only access to elements
- ◆ bidirectional iterator: supports both  $++p$  and  $-p$
- ◆ random-access iterator: supports both  $p+i$  and  $p-i$

# STL Iterators in C++

---

□ Each STL container type `C` supports iterators:

- `C::iterator` – read/write iterator type
- `C::const_iterator` – read-only iterator type
- `C.begin()`, `C.end()` – return start/end iterators

□ This iterator-based operators and methods:

- `*p`: access current element
- `++p`, `--p`: advance to next/previous element
- `C.assign(p, q)`: replace `C` with contents referenced by the iterator range `[p, q)` (from `p` up to, but not including, `q`)
- `insert(p, e)`: insert `e` prior to position `p`
- `erase(p)`: remove element at position `p`
- `erase(p, q)`: remove elements in the iterator range `[p, q)`

# Back to Iterator: STL Iterator-based Functions

---

`vector(p, q)`: Construct a vector by iterating between  $p$  and  $q$ , copying each of these elements into the new vector.

`assign(p, q)`: Delete the contents of  $V$ , and assigns its new contents by iterating between  $p$  and  $q$  and copying each of these elements into  $V$ .

`insert(p, e)`: Insert a copy of  $e$  just prior to the position given by iterator  $p$  and shifts the subsequent elements one position to the right.

`erase(p)`: Remove and destroy the element of  $V$  at the position given by  $p$  and shifts the subsequent elements one position to the left.

`erase(p, q)`: Iterate between  $p$  and  $q$ , removing and destroying all these elements and shifting subsequent elements to the left to fill the gap.

`clear()`: Delete all these elements of  $V$ .

# STL Containers and Algorithms

---

#include <algorithm>

`sort(p, q)`: Sort the elements in the range from  $p$  to  $q$  in ascending order. It is assumed that less-than operator (“<”) is defined for the base type.

`random_shuffle(p, q)`: Rearrange the elements in the range from  $p$  to  $q$  in random order.

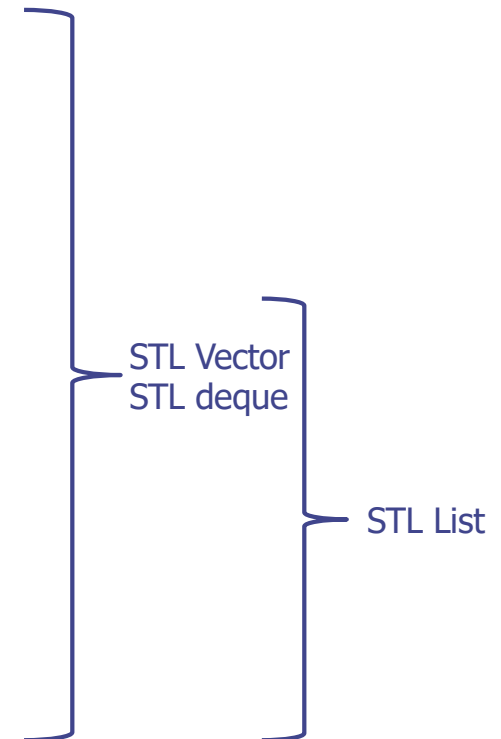
`reverse(p, q)`: Reverse the elements in the range from  $p$  to  $q$ .

`find(p, q, e)`: Return an iterator to the first element in the range from  $p$  to  $q$  that is equal to  $e$ ; if  $e$  is not found,  $q$  is returned.

`min_element(p, q)`: Return an iterator to the minimum element in the range from  $p$  to  $q$ .

`max_element(p, q)`: Return an iterator to the maximum element in the range from  $p$  to  $q$ .

`for_each(p, q, f)`: Apply the function  $f$  the elements in the range from  $p$  to  $q$ .



<http://www.cplusplus.com/reference/algorithm/>

## Example Code

---

```
#include <cstdlib> // provides EXIT_SUCCESS
#include <iostream> // I/O definitions
#include <vector> // provides vector
#include <algorithm> // for sort, random_shuffle

using namespace std; // make std:: accessible

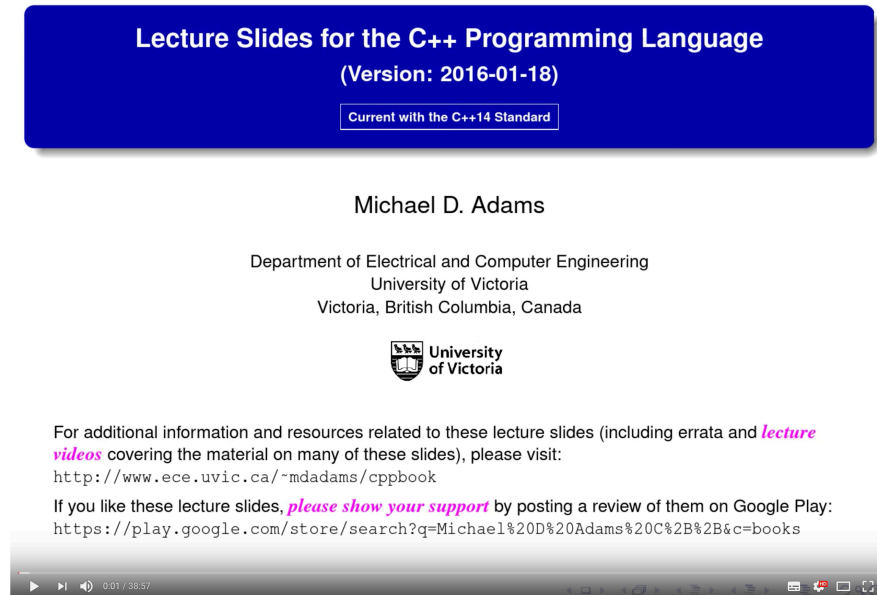
int main () {
    int a[] = {17, 12, 33, 15, 62, 45};
    vector<int> v(a, a + 6); // v: 17 12 33 15 62 45
    cout << v.size() << endl; // outputs: 6
    v.pop_back(); // v: 17 12 33 15 62
    cout << v.size() << endl; // outputs: 5
    v.push_back(19); // v: 17 12 33 15 62 19
    cout << v.front() << " " << v.back() << endl; // outputs: 17 19
    sort(v.begin(), v.begin() + 4); // v: (12 15 17 33) 62 19
    v.erase(v.end() - 4, v.end() - 2); // v: 12 15 62 19
    cout << v.size() << endl; // outputs: 4

    char b[] = {'b', 'r', 'a', 'v', 'o'};
    vector<char> w(b, b + 5); // w: b r a v o
    random_shuffle(w.begin(), w.end()); // w: o v r a b
    w.insert(w.begin(), 's'); // w: s o v r a b

    for (vector<char>::iterator p = w.begin(); p != w.end(); ++p)
        cout << *p << " "; // outputs: s o v r a b
    cout << endl;
    return EXIT_SUCCESS;
}
```


# If you want to know more about iterators,

◆ Please watch this video



Lecture Slides for the C++ Programming Language  
(Version: 2016-01-18)  
Current with the C++14 Standard

Michael D. Adams  
Department of Electrical and Computer Engineering  
University of Victoria  
Victoria, British Columbia, Canada



For additional information and resources related to these lecture slides (including errata and *lecture videos* covering the material on many of these slides), please visit:  
<http://www.ece.uvic.ca/~mdadams/cppbook>

If you like these lecture slides, *please show your support* by posting a review of them on Google Play:  
<https://play.google.com/store/search?q=Michael%20Adams%20C%2B%2B&c=books>

<https://www.youtube.com/watch?v=TxufBysSPK0>

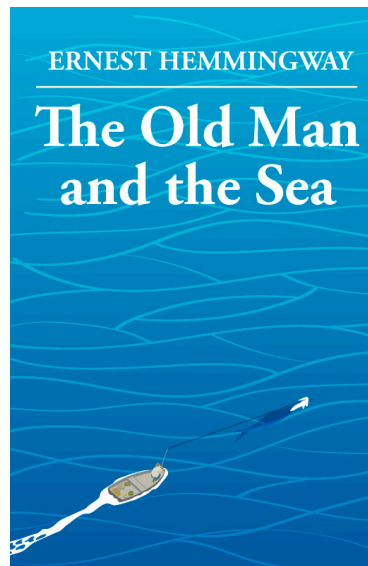
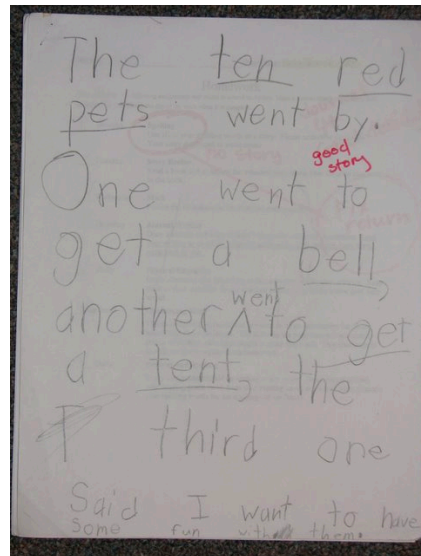
◆ Please

- I hate to answer the question “Is this included in the exam?”

# What should be your next question?

---

- ◆ Can I implement iterators in C++, in addition to just knowing how to use them?
  - Someone like the C++ STL designer
- ◆ Ch 6.2.3: Some level of explanation:
  - Beyond the topic of this class
- ◆ I will be happy to discuss this if you visit my office.



# Sequences





# Sequence ADT

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- ◆ The **Sequence** ADT is the union of the Array List and Node List ADTs
- ◆ Elements accessed by
  - Index, or
  - Position
- ◆ Generic methods:
  - **size()**, **empty()**
- ◆ ArrayList-based methods:
  - **at(i)**, **set(i, o)**, **insert(i, o)**, **erase(i)**
- ◆ List-based methods:
  - **begin()**, **end()**
  - **insertFront(o)**, **insertBack(o)**
  - **eraseFront()**, **eraseBack()**
  - **insert (p, o)**, **erase(p)**
- ◆ Bridge methods:
  - **atIndex(i)**, **indexOf(p)**

# Applications of Sequences

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- ◆ The Sequence ADT is a basic, general-purpose, data structure for storing an ordered collection of elements
- ◆ Direct applications:
  - Generic replacement for stack, queue, vector, or list
  - small database (e.g., address book)
- ◆ Indirect applications:
  - Building block of more complex data structures

Questions?